

Part I: Review of Basic Concepts

Introduction

- **Graph theory** is a powerful mathematical framework used to model and analyze complex relationships in various fields, including computer science, telecommunications, and information networks.
- In this course, we will explore the fundamental concepts and algorithms of graph theory with a focus on their practical **applications in network science**.

Applications in information networks

- **Topology Modeling:** Graphs represent network structures, nodes are routers or servers, and edges are connections.
 - *Ex: A data center network modeled as a **Fat-tree topology** ensures efficient data flow.*
- **Shortest Path Algorithms:** Used to find the most efficient route between nodes in a network.
 - *Ex: Dijkstra's algorithm helps routers compute the fastest path for data packets.*
- **Multicast & Broadcast Routing:** Ensures efficient data transmission to multiple recipients using spanning trees.
 - *Ex: IP multicast delivers a live video stream to multiple users without duplicate transmissions.*
- **Flow Algorithms:** Optimize data flow through a network by modeling capacity constraints.
 - *Ex: Ford-Fulkerson algorithm helps balance traffic across network links to avoid congestion.*

Applications in information networks

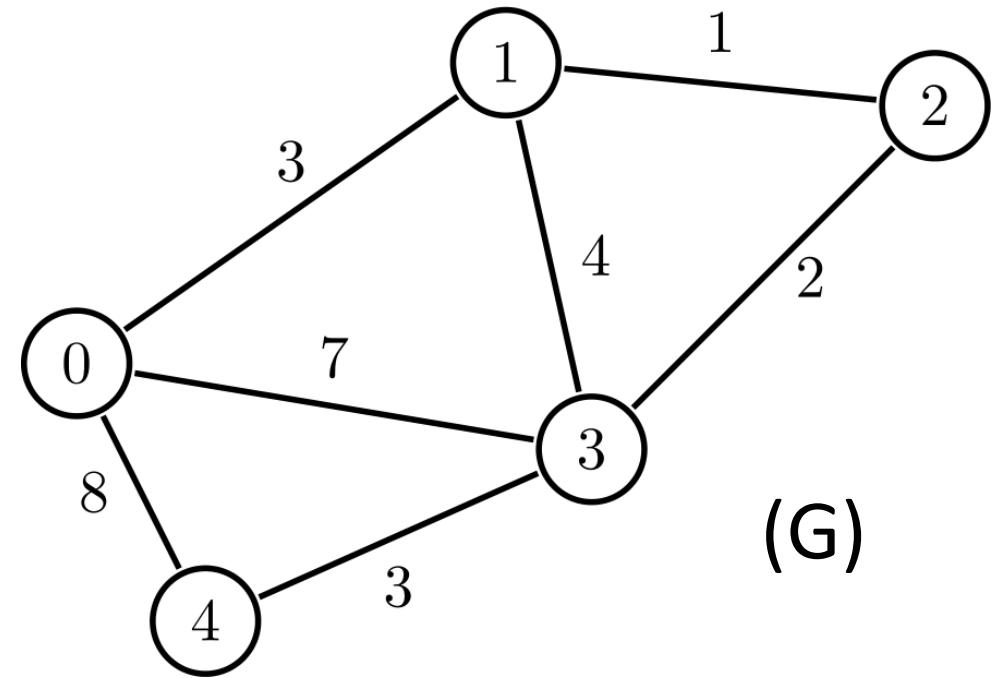
- **Connectivity & Vulnerability:** Analyzes how network failures affect connectivity and identifies weak points.
 - Ex: Finding articulation points in a network shows which nodes' failure would break connectivity.
- **Influence & Information Spread:** Models how information, trends, or viruses propagate through networks.
 - Ex: Identifying key influencers in a social media network using centrality measures.
- **Anomaly Detection:** Detects unusual patterns in network activity by monitoring graph structure changes.
 - Ex: A sudden change in node degree could indicate an attack on a server.

Graph vs Network

- Graph: A set of vertices V and edges E connecting them: $G=(V,E)$.
 - $V=\{v_1, v_2, \dots, v_n\}$, v_i is a **vertex**
 - $E=\{e_1, e_2, \dots, e_m\}$, e_j is an **edge**, $e=\{u,v\}$ is simply noted $e=uv$
 - **Adjacent** nodes have a common edge
 - **Adjacent** edges have a common node
 - If $e=uv$, the edge e and the vertices u and v are **incidents**
- Networks: Real structures that can be represented by graphs
 - Big size (nodes & links)
 - Complex structure with attributed nodes and relations

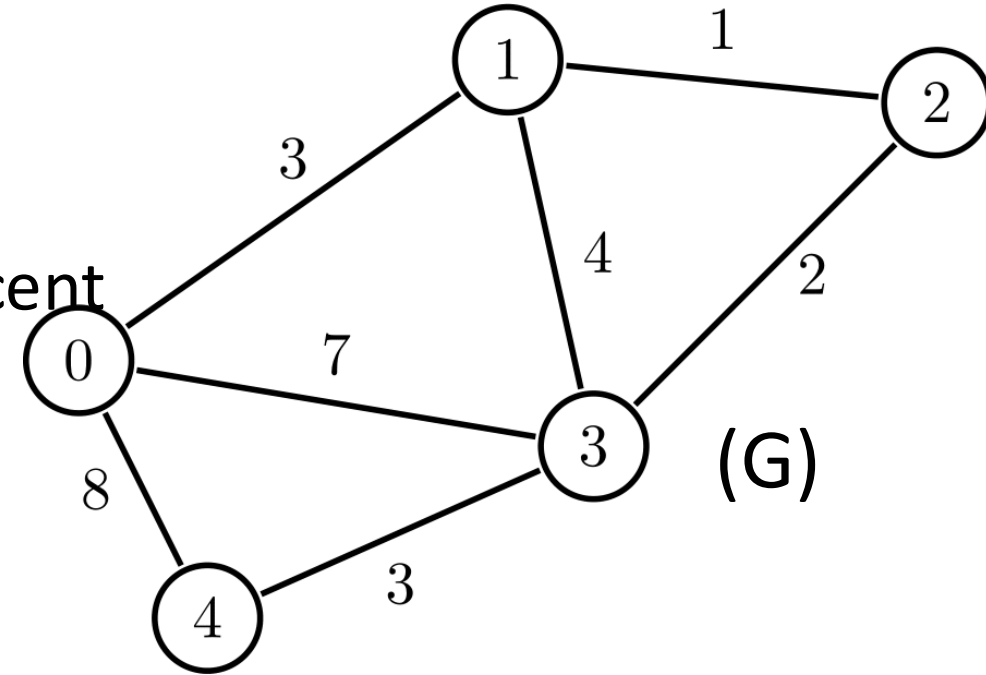
Basic metrics

- **Order:** number of vertices, $\text{Order}(G) = |V|$
 - Ex: $\text{order}(G) = 5$
- **Size:** number of edges, $\text{Size}(G) = |E|$
 - Ex: $\text{size}(G) = 7$
- **Degree:** number of edges connected to a vertex, $\text{deg}(v)$ = number of edges connected to v , $\text{deg}(G)$ = max degree of all vertices
 - $\sum_{v \in V} \text{degree}(v) = 2 \times \text{size}(G)$
 - $\sum : 3 + 2 + 4 + 3 + 2 = 14 = 2 \times 7$
- **Weight:** A value assigned to an edge
 - Ex: $\text{weight}(0,4) = 8$, $\text{weight}(1,3) = 4$



Important concepts

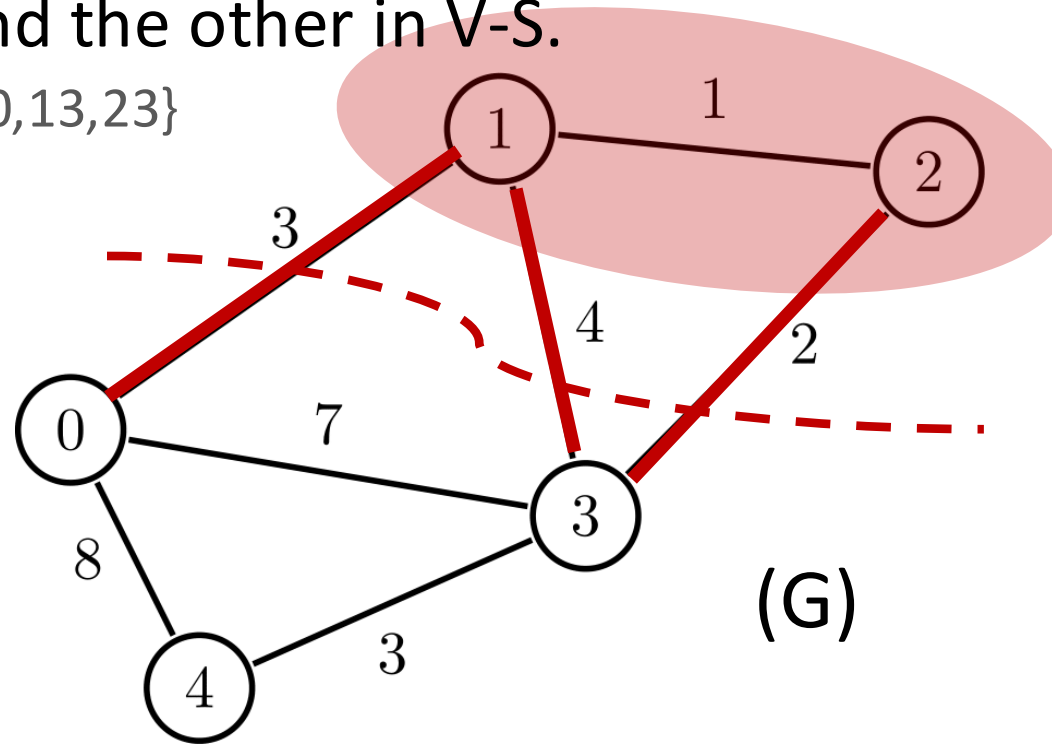
- **Walk:** A sequence of vertices pairwise adjacent
 - Ex: 0,1,3,0,1,2,3,4
 - Ex: 4,1,2,1,0 (**not a walk**)
- **Trail:** A walk with no repeated edges
 - Ex: 0,3,1,0,4
- **Path:** A trail with no repeated vertices
 - Ex: 4,3,1,2
- **Cycle:** A path that starts and ends at the same vertex
 - Ex: 2,1,0,3,2
- **Distance** between 2 vertices is the shortest path length between them
 - Ex: $\text{distance}(1,4) = 1 + 2 + 3 = 6$



Important concepts

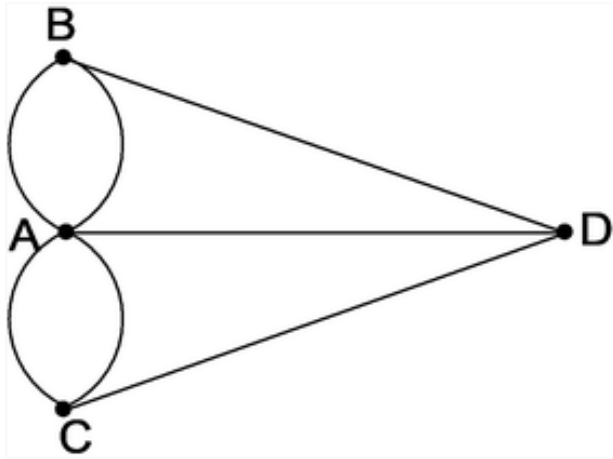
- **Co-cycle(cutset)**: set of edges whose removal increases the number of connected components in the graph.
It is often associated with a partition of the vertex set V into two disjoint subsets S and $V-S$. The cocycle consists of all edges that have one endpoint in S and the other in $V-S$.

- Ex: $\text{co-cycle}(\{1,2\}) = \{10,13,23\}$

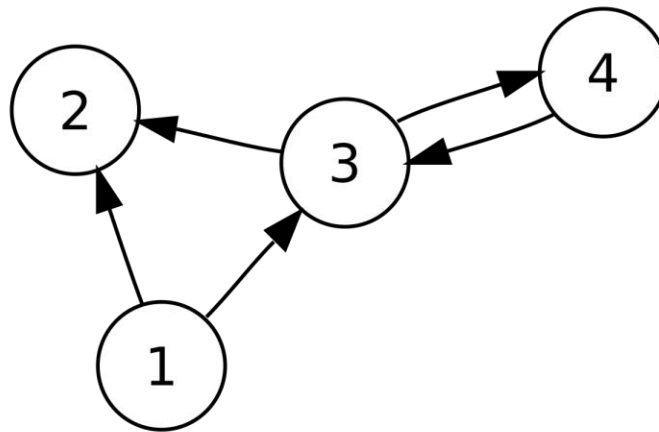


Types of graphs

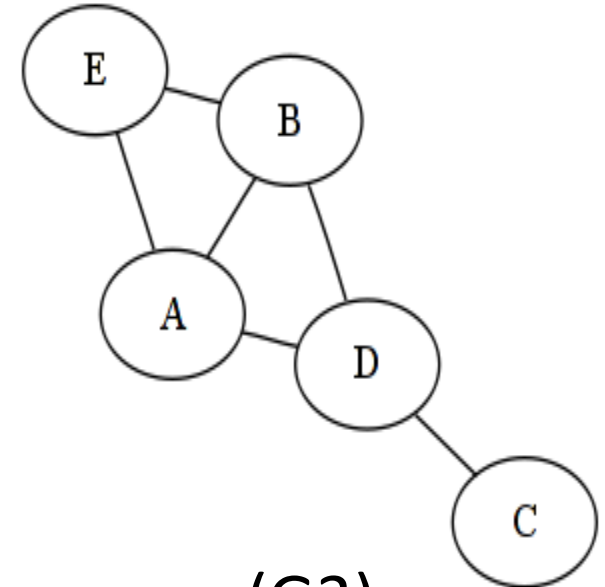
- **Indirected:** Edges have no direction (G1)
- **Directed(DiGraph):** Edges have direction (G2)
- **Multi-Graph:** Allows multiple edges between two vertices (G3)



(G1)



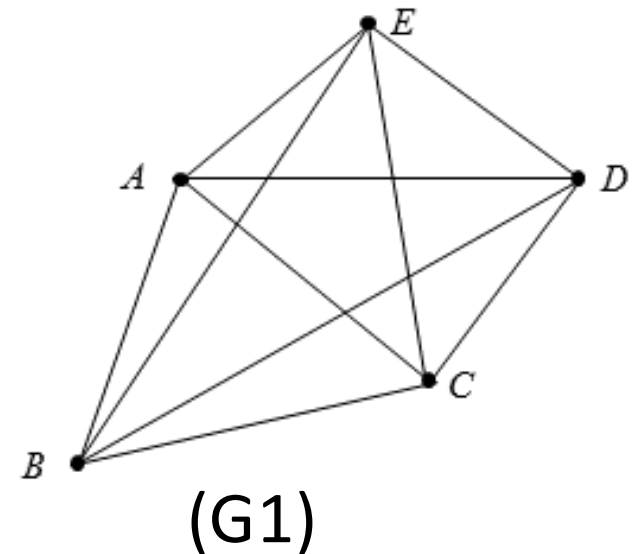
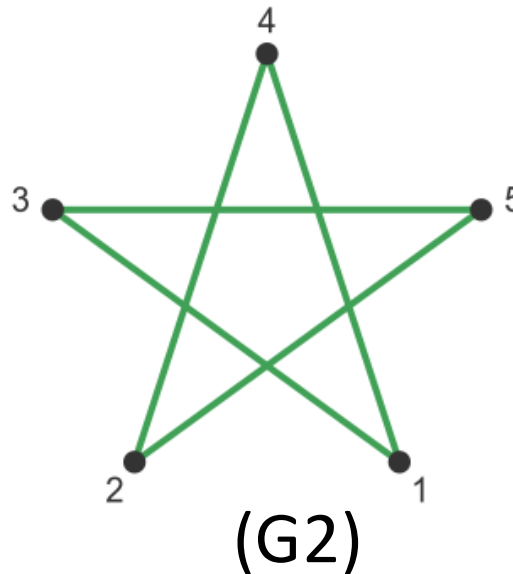
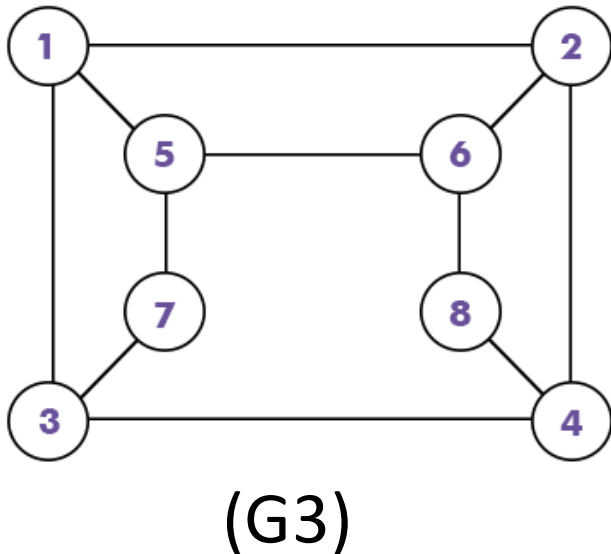
(G2)



(G3)

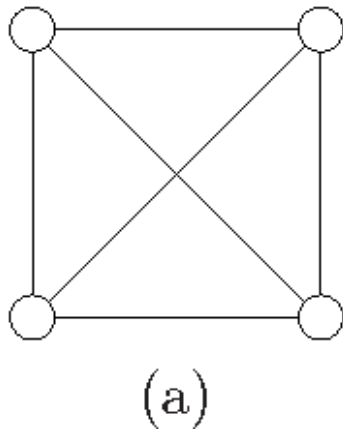
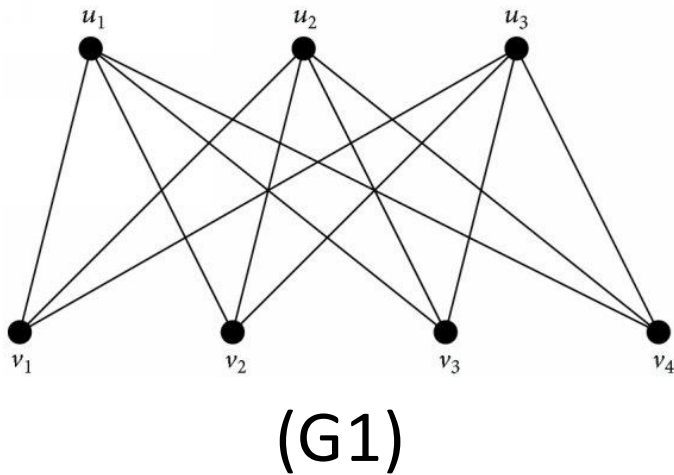
Types of graphs

- **Complete (k_n):** Every pair of vertices is connected (G1)
- **Regular(n-regular):** All vertices have the same degree (G2)
- **Bipartite:** Vertices can be divided into two disjoint sets, edges link only vertices from one set to the other set (G3)

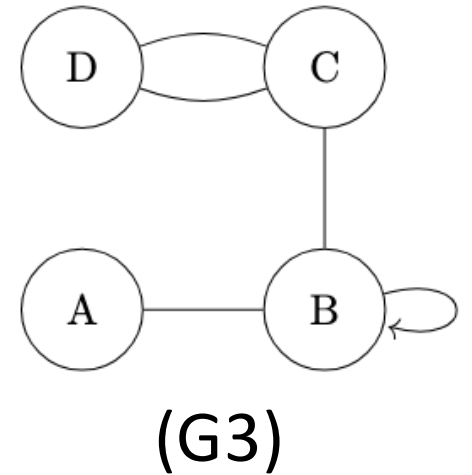
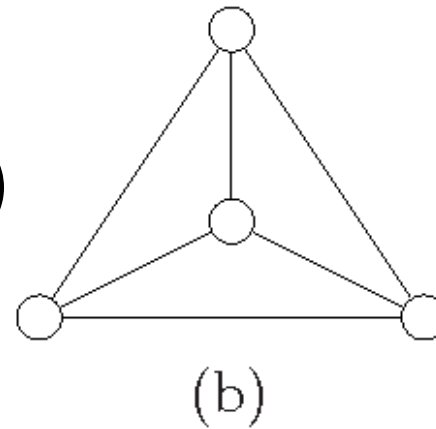


Types of graphs

- **Complete bipartite:** A bipartite graph with all possible edges (G1)
- **Planar:** Can be drawn without edge crossings (G2). Delimited areas are named **Faces**, the graph of faces is named **Dual graph**
 - $\sum \text{degree}(\text{face}) = 2 \times \text{size}(G)$
 - $\text{Order}(G) - \text{size}(G) + \# \text{ of face} = 2$
- **Simple:** No loops or multiple edges (G3 is not a simple graph)

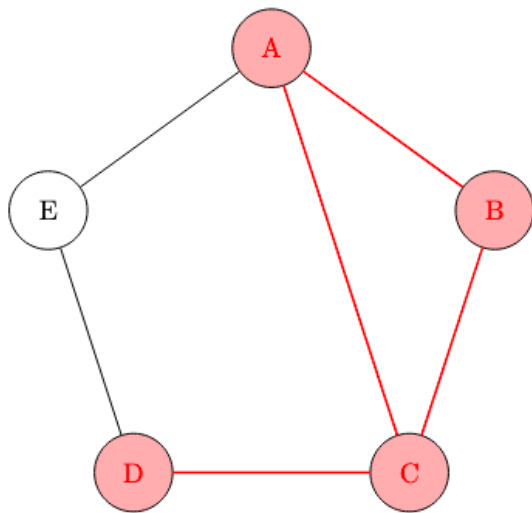


(G2)

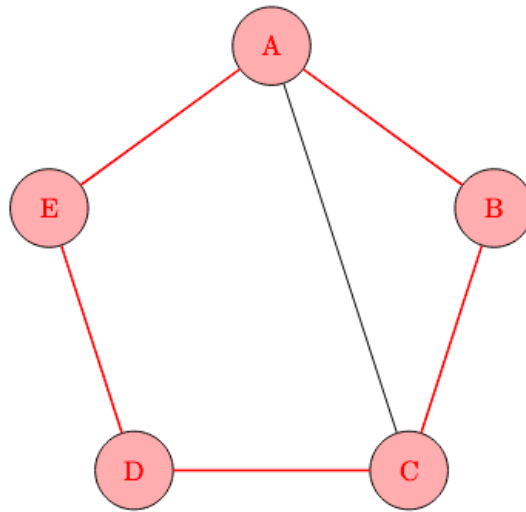


Types of graphs

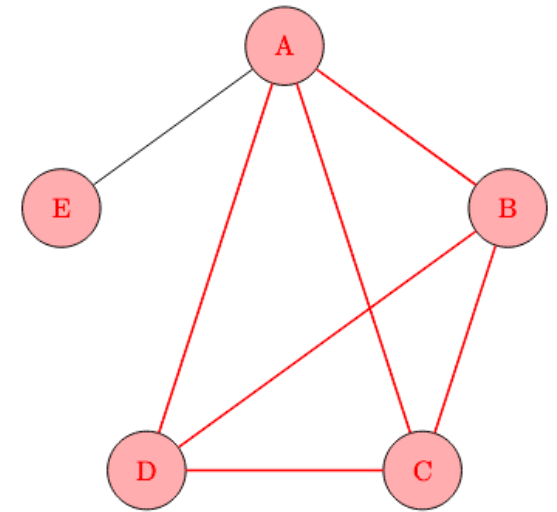
- **Subgraph:** A subset of a graph's vertices with incident edges (G1)
- **Partial:** A subgraph that retains all vertices but removes some edges (G2)
- **Clique:** A subgraph where every pair of vertices is connected (G3)



(G1)



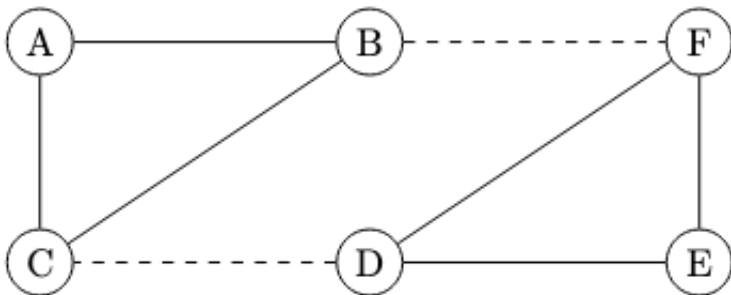
(G2)



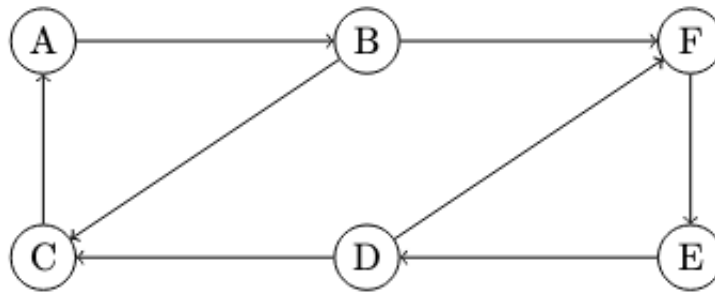
(G3)

Connectivity

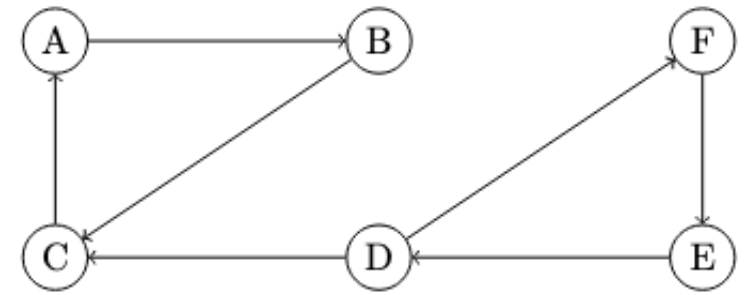
- Undirected graph: A graph is **connected** if there is a path between every pair of vertices and **disconnected** else (G1).
- Directed graph: A DiGraph can be disconnected, weakly or strongly connected.
 - A digraph is **strongly connected** if every vertex can reach every other vertex (G2).
 - A digraph is **weakly connected** if it will be connected after transforming their arcs to edges (G3).



(G1)



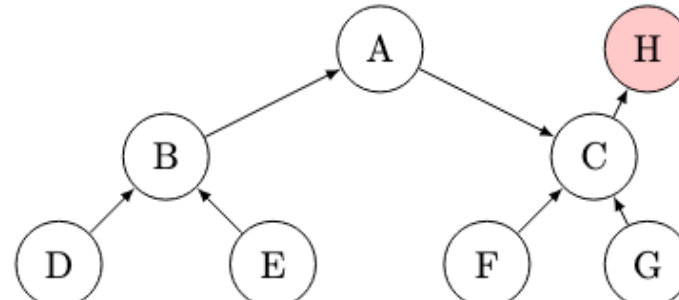
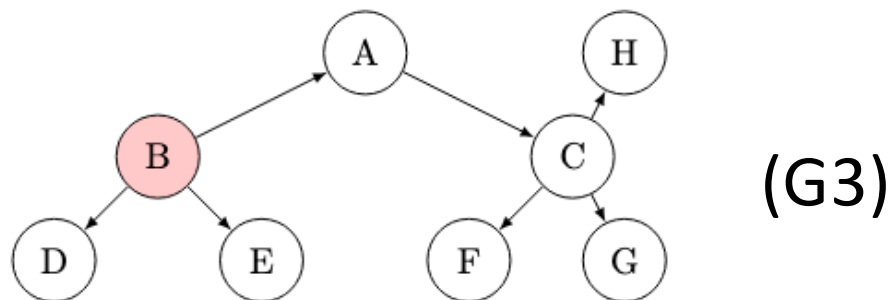
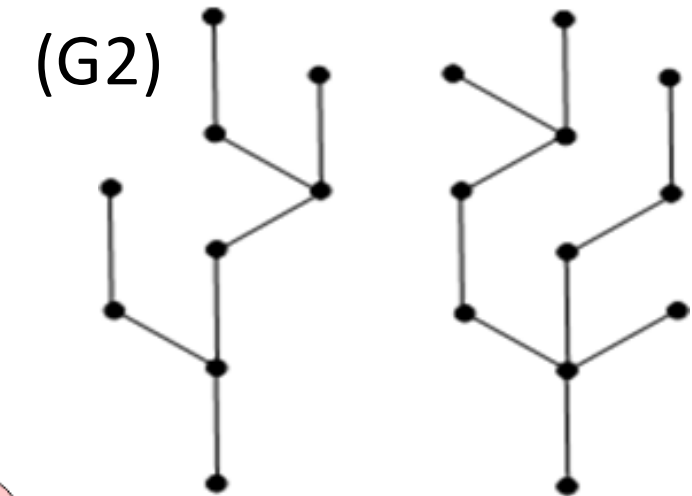
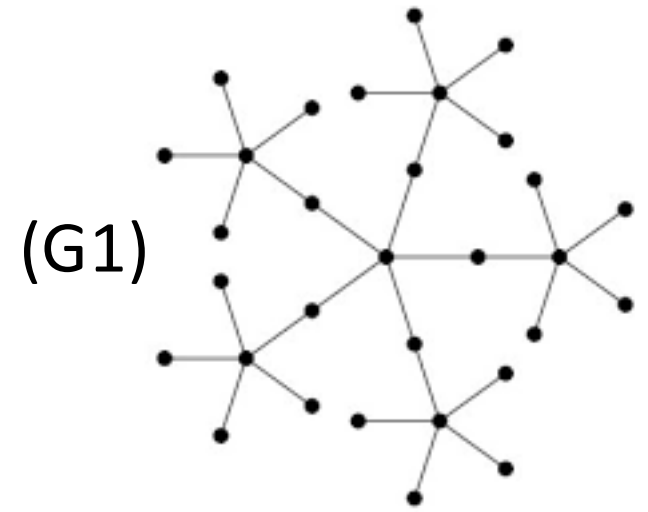
(G2)



(G3)

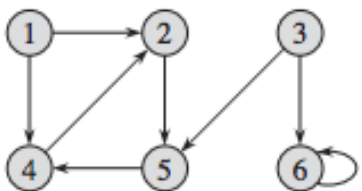
Trees

- **Tree:** A connected acyclic graph (G1)
 - Size=order-1
- **Forest:** A collection of disjoint trees (G2)
- **Arborescence/Anti-Arborescence:** A directed tree with a single root, from which all edges point outward/inward (G3)
- **Root/Anti-Root:** A vertex that allows reaching/reachable by all other vertices

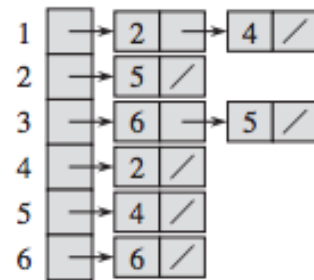


Computer representation of graphs

- **Adjacency matrix:** A $|V| \times |V|$ matrix where $A[i][j]=1$ if an edge exists between v_i and v_j .
- **Adjacency list:** A list where each vertex has a set of adjacent vertices.
- **Incidence matrix:** A $|V| \times |E|$ matrix where rows represent vertices and columns represent edges.
- **Adjacency matrix Power:** The matrix power A^k shows the number of paths of length k between vertices.



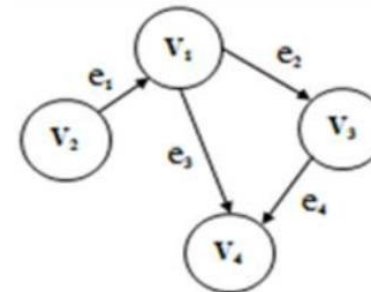
(a)



(b)

	1	2	3	4	5	6
1	0	1	0	1	0	0
2	0	0	0	0	1	0
3	0	0	0	0	1	1
4	0	1	0	0	0	0
5	0	0	0	1	0	0
6	0	0	0	0	0	1

(c)



(i)

	e_1	e_2	e_3	e_4
v_1	-1	1	1	0
v_2	1	0	0	0
v_3	0	-1	0	1
v_4	0	0	-1	-1

(ii)